# Towards a parallel domain decomposition solver for immersed boundary finite element method 

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## Outline

Introduction

Parallel adaptive mesh refinement

Multilevel BDDC method

Immersed boundary FEM

Numerical results

Conclusion and outlooks

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## Parallel adaptive mesh refinement

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## Conclusion and outlooks

1 Adaptivity and higher order finite elements
[P. Kůs, P. Šolín, D. Andrš, Arbitrary-level hanging nodes for adaptive hp-FEM approximations in 3D, JCAM, 270, pp. 121-133, 2014.]


2 Nonoverlapping domain decomposition and parallel computing
[B. Sousedík, J. Šístek, and J. Mandel, Adaptive-Multilevel BDDC and its parallel implementation, Computing, 95 (12), pp. 1087-1119, 2013.]


3 Immersed boundary FEM

[T. Rüberg, F. Cirak, and J.M. Garcia-Aznar, An unstructured immersed finite element method for nonlinear solid mechanics, Advanced Modeling and Simulation in Engineering Sciences, 2016.]

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C. Burstedde, L. Wilcox, and O. Ghattas, p4est: Scalable Algorithms for Parallel Adaptive Mesh Refinement on Forests of Octrees, SIAM J. Sci. Comput., 3 (33), pp. 1103-1133, 2011.


1 Hanging nodes

- Hanging nodes have to be eliminated
- They can also appear at the subdomain interface

2 Shape of the subdomains

- Subdomains might be disconnected or only loosely coupled (e.g. by one node in elasticity)


## Adding rules to the game

## Assumptions on the mesh

- only level-1 hanging nodes allowed
- 2:1 rule
- equal order shape functions, i.e. no $h p$, but higher $p$ fine

O.K.

not O.K.


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## An abstract problem

$$
u \in U: a(u, v)=\langle f, v\rangle \quad \forall v \in U
$$

- $a(\cdot, \cdot)$ symmetric positive definite form on $U$
- $\langle\cdot, \cdot\rangle$ is inner product on $U$
- $U$ is finite dimensional space (typically finite element functions)
- A symmetric positive definite matrix on $U$
- A large, sparse, condition number $\kappa(A)=\frac{\lambda_{n}}{\lambda^{n}}=O\left(1 / h^{2}\right)$


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- $U$ is finite dimensional space (typically finite element functions)


## Matrix form

$$
u \in U: A u=f
$$

- $A$ symmetric positive definite matrix on $U$
- $A$ large, sparse, condition number $\kappa(A)=\frac{\lambda_{\text {max }}}{\lambda_{\text {min }}}=\mathcal{O}\left(1 / h^{2}\right)$
- idea goes back to substructuring - a trick used in seventies to fit larger FE problems into memory

- $\Omega_{1}, \Omega_{2} \ldots$ subdomains (substructures)
- $\Gamma \ldots$ interface
- unknowns at interface are shared by more subdomains, remaining (interior) unknowns belong to a single subdomain
- the first step is reduction of the problem to the interface I
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- $\Omega_{1}, \Omega_{2} \ldots$ subdomains (substructures)
- $\Gamma$. . . interface

■ unknowns at interface are shared by more subdomains, remaining (interior) unknowns belong to a single subdomain

- the first step is reduction of the problem to the interface $\Gamma$

Formation of the interface problem

- recall the matrix problem

$$
A u=f
$$

- reorder unknowns so that those at interior $u_{o}^{1}$ and $u_{o}^{2}$ are first, then interface $u_{\Gamma}$

- eliminate interior unknowns - subdomain by subdomain $=$ in parallel


> assembly
assembly
assembly
assembly

- recall the matrix problem

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A u=f
$$

- reorder unknowns so that those at interior $u_{o}^{1}$ and $u_{o}^{2}$ are first, then interface $u_{\Gamma}$

$$
\left[\begin{array}{ccc}
A_{o o}^{1} & & A_{o \Gamma}^{1} \\
A_{\Gamma o}^{1} & A_{o o}^{2} & A_{\sigma \Gamma}^{2} \\
A_{\Gamma \Gamma}^{2} & A_{\Gamma \Gamma}
\end{array}\right]\left[\begin{array}{l}
u_{o}^{1} \\
u_{o}^{2} \\
u_{\Gamma}
\end{array}\right]=\left[\begin{array}{c}
f_{o}^{1} \\
f_{o}^{2} \\
f_{\Gamma}
\end{array}\right]
$$

- eliminate interior unknowns - subdomain by subdomain $=$ in parallel
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A_{\Gamma o}^{1} & A_{\Gamma o}^{2} & A_{\Gamma \Gamma}
\end{array}\right]\left[\begin{array}{c}
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f_{o}^{2} \\
f_{\Gamma}
\end{array}\right]
$$

- eliminate interior unknowns - subdomain by subdomain $=$ in parallel

$$
\begin{aligned}
& {\left[\begin{array}{lll}
A_{o o}^{1} & & A_{o \Gamma}^{1} \\
& A_{o o}^{2} & A_{o \Gamma}^{2} \\
& & S
\end{array}\right]\left[\begin{array}{c}
u_{o}^{1} \\
u_{o}^{2} \\
u_{\Gamma}
\end{array}\right]=\left[\begin{array}{c}
f_{o}^{1} \\
f_{o}^{2} \\
g
\end{array}\right]} \\
& S=\sum_{\text {assembly }} A_{\Gamma \Gamma}^{i}-A_{\Gamma o}^{i}\left(A_{o o}^{i}\right)^{-1} A_{o \Gamma}^{i}=\sum_{\text {assembly }} S^{i} \\
& g=\sum_{\text {assembly }} f_{\Gamma}^{i}-A_{\Gamma o}^{i}\left(A_{o o}^{i}\right)^{-1} f_{o}^{i}=\sum_{\text {assembly }} g^{i}
\end{aligned}
$$



- interface $\Gamma$

Reduced (Schur complement) problem on interface $\Gamma$

- solved by PCG

- interface $\Gamma$


## Reduced (Schur complement) problem on interface $\Gamma$

$$
\begin{gathered}
S u_{\Gamma}=g \\
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\end{gathered}
$$

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\end{gathered}
$$

- solved by PCG


## A practical algorithm of iterative substructuring

- In setup:

1 factorize matrix $A_{o o}$ (block diagonal $=$ in parallel)
$\sqrt{2}$ form condensed right-hand side by solving

$$
A_{o o} h=f_{o},
$$

and inserting $g=f_{\Gamma}-A_{\Gamma \circ} h$.

- In each iteration, for given $p$ construct $S p$ as


1 Solve (in parallel) discrete Dirichlet problem

2 Get $S p$ (in parallel) as

- After iterations, for given $u_{\Gamma}$, resolve (in parallel) interior unknowns by back-substitution in


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$$
\left[\begin{array}{ll}
A_{o o} & A_{o \Gamma} \\
A_{Г o} & A_{\Gamma \Gamma}
\end{array}\right]\left[\begin{array}{l}
w \\
p
\end{array}\right]=\left[\begin{array}{c}
0 \\
S p
\end{array}\right]
$$

1 Solve (in parallel) discrete Dirichlet problem

$$
A_{o o} w=-A_{o \Gamma} p
$$

$\sqrt{2}$ Get $S p$ (in parallel) as

$$
S p=A_{\Gamma o} w+A_{\Gamma \Gamma} p
$$

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$$
A_{o o} u_{o}=f_{o}-A_{o \Gamma} u_{\Gamma}
$$

## The BDDC preconditioner



- Balancing Domain Decomposition based on Constraints [Dohrmann (2003)], [Cros (2003)], [Fragakis, Papadrakakis (2003)]
- continuity at corners, and of averages (arithmetic or weighted) over edges or faces considered
- enough constraints to fix floating subdomains - $a(\cdot, \cdot)$ symmetric positive definite on $\widetilde{W}$


## The BDDC preconditioner



continuous at all nodes at interface
 $\subset$


no continuity at interface

- Balancing Domain Decomposition based on Constraints [Dohrmann (2003)], [Cros (2003)], [Fragakis, Papadrakakis (2003)]
- continuity at corners, and of averages (arithmetic or weighted) over edges or faces considered
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Variational form of $M_{B D D C}: r \longmapsto u$

$$
\begin{aligned}
w \in \widetilde{W}: \quad a(w, z) & =\langle r, E z\rangle \quad \forall z \in \widetilde{W} \\
u & =E w
\end{aligned}
$$

## Local energy minimization problems

On each subdomain - coarse degrees of freedom - basis functions $\Psi^{i}$ prescribed values of coarse degrees of freedom, minimal energy elsewhere,

$$
\left[\begin{array}{cc}
A^{i} & C^{i T} \\
C^{i} & 0
\end{array}\right]\left[\begin{array}{l}
\Psi^{i} \\
\Lambda^{i}
\end{array}\right]=\left[\begin{array}{l}
0 \\
I
\end{array}\right] .
$$

- $A^{i}$... local subdomain stiffness matrix
- $C^{i} \ldots$ matrix of constraints - selects unknowns into coarse degrees of freedom Matrix of coarse problem $A_{C}$ assembled from local matrices $A_{C i}=\Psi^{i T} A^{i} \Psi^{i}=-\Lambda^{i}$.

coarse basis fun.


## Assumption on solvability of local saddle-point systems

$$
\text { null } A^{i} \cap \operatorname{null} C^{i}=\{\mathbf{0}\}
$$

- invertibility of

$$
\left[\begin{array}{cc}
A^{i} & C^{i T} \\
C^{i} & 0
\end{array}\right]
$$

follows, see e.g. [Benzi, Golub, Liesen (2005)]

- satisfied if enough constraints on continuity in $\widetilde{W}$ are selected


## Disconnected and loosely coupled subdomains

- more constrains needed in $C^{i}$
- more coarse basis functions $\Psi^{i}$

- local nullspaces more complicated
- more constrains needed for each subdomain
- detect graph components of subdomain mesh
P. Kůs, J. Šístek. Coupling parallel adaptive mesh refinement with a nonoverlapping domain decomposition solver. Advances in Engineering Software, 110:34-54, 2017.

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Get residual at interface nodes $r_{\Gamma}^{(k)}=g-S u_{\Gamma}^{(k)}$ and produce preconditioned residual $z_{\Gamma}^{(k)}=M_{B D D C} r_{\Gamma}^{(k)}$ by

1 1. Distribution of residual
subdomain problems (global) coarse problem

2 Correction of solution


3 Combination of subdomain and coarse corrections


## One step of BDDC

Get residual at interface nodes $r_{\Gamma}^{(k)}=g-S u_{\Gamma}^{(k)}$ and produce preconditioned residual $z_{\Gamma}^{(k)}=M_{B D D C} r_{\Gamma}^{(k)}$ by
11 . Distribution of residual
subdomain problems (global) coarse problem

$$
\begin{gathered}
\text { for } i=1, \ldots, N \\
r^{i}=E^{i T} r_{\Gamma}^{(k)}
\end{gathered}
$$

$$
r_{C}=\sum_{i=1}^{N} R_{C}^{i T} \Psi^{i * T} E^{i T} r_{\Gamma}^{(k)}
$$

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$$
\left[\begin{array}{cc}
A^{i} & C^{i T} \\
C^{i} & 0
\end{array}\right]\left[\begin{array}{c}
z^{i} \\
\mu^{i}
\end{array}\right]=\left[\begin{array}{c}
r^{i} \\
0
\end{array}\right] \quad A_{C} u_{C}=r_{C}
$$

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3 Combination of subdomain and coarse corrections

$$
z_{\Gamma}^{(k)}=\sum_{i=1}^{N} E^{i}\left(\Psi^{i} R_{C}^{i} u_{C}+z^{i}\right)
$$

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- an approach to avoiding mesh generation in FEM simulations
- robustness for geometric resolution through implicit description of the boundary

Nonstandard features compared to FEM

- weak enforcement of Dirichiet boundary conditions
- nonstandard quadrature rules
- ill-conditioning of stiffness matrices due to cut cells

T. Rüberg, F. Cirak, and J.M. Garcia-Aznar, An unstructured immersed finite element method for nonlinear solid mechanics, Advanced Modeling and Simulation in Engineering Sciences, 2016.
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$$
u \in H^{1}(\Omega): \quad a(u, v)=l(v) \quad \forall v \in H^{1}(\Omega)
$$

## Penalty method

$$
\begin{aligned}
a(u, v) & =\int_{\Omega} \nabla u \cdot \nabla v \mathrm{~d} \Omega+\int_{\Gamma_{D}} \gamma u v \mathrm{~d} \Gamma \\
l(v) & =\int_{\Omega} f v \mathrm{~d} \Omega+\int_{\Gamma_{N}} \bar{t} v \mathrm{~d} \Gamma+\int_{\Gamma_{D}} \gamma \bar{u} v \mathrm{~d} \Gamma
\end{aligned}
$$

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\end{aligned}
$$

Nitsche method

$$
\begin{aligned}
a(u, v) & =\int_{\Omega} \nabla u \cdot \nabla v \mathrm{~d} \Omega-\int_{\Gamma_{D}}(\mathbf{n} \cdot \nabla u) v \mathrm{~d} \Gamma \\
& -\int_{\Gamma_{D}}(\mathbf{n} \cdot \nabla v) u \mathrm{~d} \Gamma+\int_{\Gamma_{D}} \gamma u v \mathrm{~d} \Gamma \\
l(v) & =\int_{\Omega} f v \mathrm{~d} \Omega+\int_{\Gamma_{N}} \bar{t} v \mathrm{~d} \Gamma-\int_{\Gamma_{D}}(\mathbf{n} \cdot \nabla v) \bar{u} \mathrm{~d} \Gamma+\int_{\Gamma_{D}} \gamma \bar{u} v \mathrm{~d} \Gamma
\end{aligned}
$$

- $\gamma=\gamma_{0} \alpha / h$


## Parallel FEM solver with AMR and embedded domains

- experimental in-house code
- implicit geometry description
- C++ with MPI


## p4est mesh manager for AMR

- rebalancing based on Z-curves
- ANSI C + MPI
- open-source (GPL)
- scalability reported for $1 e 5-1$ e6 cores
http://www.p4est.org


## BDDCML equation solver

- Adaptive-Multilevel BDDC
- Fortran 95 + MPI
- open-source (LGPL)
- current version 2.5 (8/6/'15)
- tested on up to 65 e 3 cores and 2 e 9 unknowns
http://www.math.cas.cz/~sistek/ software/bddcml.html


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## Adaptive mesh refinement

## Internal layer benchmark

$$
\begin{aligned}
-\Delta u & =f \quad \text { on } \quad(0,1)^{d} \\
u & =\arctan \left(s \cdot\left(r-\frac{\pi}{3}\right)\right)
\end{aligned}
$$

■ solution exhibits sharp internal layer

- $r$ is a distance from a given point

■ $s$ controls "steepness" of the layer


## Adaptivity in 3D on 8 subdomains

- Adaptivity tested for element orders 1-4 (showed order 1)

■ Guided by exact solution, using $H^{1}$ semi-norm for error calculation


Iteration 3, mesh and solution

## Adaptivity in 3D on 8 subdomains

- Adaptivity tested for element orders 1-4 (showed order 1)

■ Guided by exact solution, using $H^{1}$ semi-norm for error calculation


Iteration 5, mesh and solution

## Adaptivity in 3D on 8 subdomains

- Adaptivity tested for element orders 1-4 (showed order 1)

■ Guided by exact solution, using $H^{1}$ semi-norm for error calculation


Iteration 8, mesh and solution

## Adaptivity in 3D, 3-level BDDC, linear elements

## Adaptive mesh refinement

| subs. | size | loc. size | PCG its. | time set-up [s] | time PCG [s] |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $2048 / 46$ | 4913 | 2 | $\mathbf{9}$ | 2.8 | 1.1 |
| $2048 / 46$ | 8594 | 4 | 29 | 0.55 | 1.6 |
| $2048 / 46$ | $2.5 \cdot 10^{4}$ | 12 | 53 | 0.6 | 3 |
| $2048 / 46$ | $1.3 \cdot 10^{5}$ | 63 | $\mathbf{6 0}$ | 0.67 | 3.5 |
| $2048 / 46$ | $7.0 \cdot 10^{5}$ | 342 | 54 | 0.89 | 3.5 |
| $2048 / 46$ | $3.0 \cdot 10^{6}$ | 1445 | 56 | 1.6 | 4.8 |
| $2048 / 46$ | $1.4 \cdot 10^{7}$ | 6623 | 55 | 2.9 | 10 |
| $2048 / 46$ | $6.4 \cdot 10^{7}$ | $3.1 \cdot 10^{4}$ | 55 | 10 | 33 |
| $2048 / 46$ | $2.9 \cdot 10^{8}$ | $1.4 \cdot 10^{5}$ | 56 | 61 | 130 |
| $2048 / 46$ | $1.3 \cdot 10^{9}$ | $6.3 \cdot 10^{5}$ | $\mathbf{5 1}$ | 565 | 521 |

■ run on Salomon@IT4I
P. Kůs, J. Šístek. Coupling parallel adaptive mesh refinement with a nonoverlapping domain decomposition solver. Advances in Engineering Software, 110:34-54, 2017.


Illustrative mesh, 100K DOFs, 20 subdomains, shown 1 - 20




Illustrative mesh, 100K DOFs, 20 subdomains, shown 1-18




Illustrative mesh, 100K DOFs, 20 subdomains, shown 1-16




Illustrative mesh, 100K DOFs, 20 subdomains, shown 1-14




Illustrative mesh, 100K DOFs, 20 subdomains, shown 1 - 12




Illustrative mesh, 100K DOFs, 20 subdomains, shown 1 - 10




Illustrative mesh, 100K DOFs, 20 subdomains, shown 1-8.

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## Our use of direct solver

## Two matrix factorizations

- $50 \%$ runtime (of the solver itself) spent in matrix factorizations, another $30 \%$ in back substitions (in iterations).
- We factorize two matrices on each subdomain. First $A_{o o}$ with inner DOFs, the other with approx. 20\% extra rows and columns, a saddle point problem. The first matrix is a submatrix of the second one.
- For Poisson problem and linear elasticity we use Cholesky and $L D L^{T}$ decompositions, respectively. LU with pivoting used for non-symmetric problems.

$$
\left[\begin{array}{ccc}
A_{o o} & A_{o \Gamma} & 0 \\
A_{\Gamma o} & A_{\Gamma \Gamma} & C_{\Gamma}^{T} \\
0 & C_{\Gamma} & 0
\end{array}\right]
$$

- We need to invert $A_{o o}$ and the 3 by 3 block matrix on each subdom.
- For eliptic problems, $A_{o o}$ is positive definite, the first 2 by 2 block matrix is positive semidefinite and the whole 3 by 3 block matrix is a saddle point problem.
- Only the part of $C_{i}$ corresponding to the interface $\Gamma$ is nonzero $\left(C_{\Gamma}\right)$.


## Our use of direct solver

## Distibution of subdomains

- Currently we run pure MPI with one subdomain per MPI rank, trying to balance the subdomain sizes.
- This is sometimes difficult to achieve when using the parallel adaptivity of the mesh and many subdomain components.
- Having variable number of subdomains per node is appealing from various reasons, but would require a major code refactoring.


## Thank you for your attention!

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## BDDCML library

http://users.math.cas.cz/~sistek/software/bddcml.html

